

# Lucente Stabile Atkins

## Encryption Algorithm

May 04, 2020

1 Gauss'  
Generalization of  
Wilson's Theorem

2 Key Exchange

3 Choosing  $U(n)$

4 Encryption

5 Decryption

# Gauss' Generalization of Wilson's Theorem

Wilson proved that, if  $p$  is prime:

$$(p - 1)! \equiv -1 \pmod{p}$$

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# Gauss' Generalization of Wilson's Theorem

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Wilson proved that, if  $p$  is prime:

$$(p - 1)! \equiv -1 \pmod{p}$$

Gauss then generalized, if  $U(n)$  is cyclic and  $\alpha \in U(n)$ :

$$\alpha_1 \cdot \alpha_2 \cdot 3 \dots \alpha_{\phi(n)-1} \equiv -1 \pmod{n}$$

# Key Exchange

Diffie Hellman  
Quantum key exchange  
Spies in a field

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# Key Exchange

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Diffie Hellman  
Quantum key exchange  
Spies in a field

Or a company can keep their chosen group a secret which means that no key exchange is needed.

# Key Exchange

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In our example:

$$k = 53$$

# Choosing $U(n)$

Since  $k = 53$ , and  $n = p^t$  or  $n = 2p^t$ ,

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# Choosing $U(n)$

Since  $k = 53$ , and  $n = p^t$  or  $n = 2p^t$ ,

Check:

$$54 = 2 \cdot 3^3 = 2 \cdot p^t$$

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# Choosing $U(n)$

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So we choose our group to be  $U(54)$ . Recall that the elements of our group are  $\alpha_i \in \mathbf{N}$  that are relatively prime with  $n$ :

$\{1, 5, 7, 11, 12, 17, 19, 23, 25, 29, 31, 35, 37, 41, 43, 47, 49, 53\}$

# Encryption

Encrypting the letter G:

Since G is the seventh letter in the alphabet we map G to the seventh element of the group  $U(54)$ . So  $C \mapsto 19$ . The sender first multiplies each element up to 7 to obtain the ciphertext:

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# Encryption

Encrypting the letter G:

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$$1 \cdot 5 \cdot 7 \cdot 11 \cdot 13 \cdot 17 \cdot 19 = 1,616,651 \equiv 17 \pmod{54}$$

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# Encryption

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17 is the first element of the ciphertext.

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Now the sender finds the number of times we find  $-1 \pmod n$ , called  $\Sigma$ , by performing the following calculations:

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Now the sender finds the number of times we find  $-1 \pmod n$ , called  $\Sigma$ , by performing the following calculations:

$$17 \cdot 53 = 901 \pmod{54} \equiv 37$$

$$37 \cdot 49 = 1,813 \pmod{54} \equiv 31$$

$$31 \cdot 47 = 1,457 \pmod{54} \equiv -1 \rightarrow (\text{first } \Sigma)$$

$$53 \cdot 43 = 2,279 \pmod{54} \equiv 11$$

$$11 \cdot 41 = 451 \pmod{54} \equiv 19$$

$$19 \cdot 37 = 703 \pmod{54} \equiv 1$$

$$1 \cdot 35 = 35 \pmod{54} \equiv 35$$

$$35 \cdot 31 = 1,085 \pmod{54} \equiv 5$$

$$5 \cdot 29 = 145 \pmod{54} \equiv 37$$

$$37 \cdot 25 = 925 \pmod{54} \equiv 7$$

$$7 \cdot 23 = 161 \pmod{54} \equiv -1 \rightarrow (\text{second } \Sigma)$$

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So  $\Sigma = 1$

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# Frame Title

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So our ciphertext is  $C = (17, 1)$ .



# Decryption

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The receiver decrypts by starting at the largest group element and multiplying in descending order until he finds  $\Sigma + 1$  elements congruent to  $-1 \pmod n$ .

Decrypting  $C = (17, 1)$ :

$$17 \cdot 53 = 901 \quad \text{mod } 54 \equiv 37$$

$$37 \cdot 49 = 1,813 \quad \text{mod } 54 \equiv 31$$

$$31 \cdot 47 = 1,457 \quad \text{mod } 54 \equiv -1 \rightarrow (\text{first } \Sigma)$$

$$53 \cdot 43 = 2,279 \quad \text{mod } 54 \equiv 11$$

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# Decrypt

Now the receiver know that every element of  $U(n)$  has been multiplied together.

$\{1, 5, 7, 11, 12, 17, 19\}$   $[23, 25, 29, 31, 35, 37, 41, 43, 47, 49, 53]$   
Sender  $\rightarrow$   $\leftarrow$  Receiver

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$$\alpha \cdot \alpha^1 \cdot \alpha^2 \cdots \alpha^{\phi(n)-1} \equiv -1 \pmod{n}$$

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